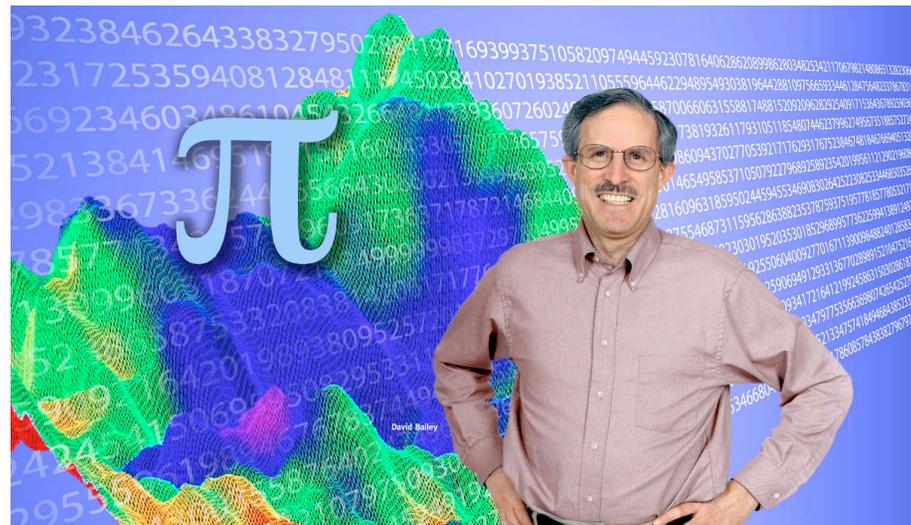


Normal Numbers

David H Bailey

Lawrence Berkeley Laboratory

<http://crd.lbl.gov/~dhbailey>



Normal numbers



Given an integer $b > 1$, a real number x is **b -normal** (or “normal base b ”) if every m -long string of digits in the base- b expansion of x appears with limiting frequency b^{-m} .

Using measure theory, it is easy to show that almost all reals are b -normal. In fact, almost all reals are b -normal for all integer bases $b > 1$.

These are widely believed to be b -normal, for all integer bases $b > 1$:

$$\pi = 3.1415926535\dots$$

$$e = 2.7182818284\dots$$

$$\text{sqrt}(2) = 1.4142135623\dots$$

$$\log(2) = 0.6931471805\dots$$

Every irrational algebraic number.

But there are no normality proofs for any of these constants, not for any base b , nor are there any non-normality results.

Until recently, normality proofs were known only for contrived examples such as Champernowne’s constant = 0.123456789101112131415... and equivalents in other bases.

A recent result for algebraic numbers



If x is algebraic of degree $d > 1$, then its binary expansion through position n must have at least $C n^{1/d}$ 1-bits, for all sufficiently large n and some C that depends on x .

Example: The first n binary digits of $\sqrt{2}$ must have at least \sqrt{n} 1-bits. In this case, the proof is easy – it follows by noting that the 1-bit count of the product of two integers is less than or equal to the product of the 1-bit counts of the two integers.

A number of other related results are established in the paper below. These results are still a far cry from full normality.

DHB, J. M. Borwein, R. E. Crandall and C. Pomerance, "On the Binary Expansions of Algebraic Numbers," *Journal of Number Theory Bordeaux*, vol. 16 (2004), pg. 487-518.

The Borwein-Plouffe observation



In 1996, Peter Borwein and Simon Plouffe of SFU in Canada observed that the following well-known formula for $\log 2$

$$\begin{aligned}\log 2 &= \sum_{n=1}^{\infty} \frac{1}{n2^n} = 0.6931471805599453094172321214581765680755 \dots_{10} \\ &= 0.101100010111001000010111111101111101000111001111011 \dots_2\end{aligned}$$

leads to a simple scheme for computing binary digits of $\log 2$ at an arbitrary starting position (here $\{ \}$ denotes fractional part):

$$\begin{aligned}\{2^d \log 2\} &= \left\{ \sum_{n=1}^d \frac{2^{d-n}}{n} \right\} + \sum_{n=d+1}^{\infty} \frac{2^{d-n}}{n} \\ &= \left\{ \sum_{n=1}^d \frac{2^{d-n} \bmod n}{n} \right\} + \sum_{n=d+1}^{\infty} \frac{2^{d-n}}{n}\end{aligned}$$

The numerator in the first portion of the RHS can be computed very rapidly using the binary algorithm for exponentiation mod n .

Fast exponentiation mod n



Problem:

What is $3^{17} \bmod 10$?

Algorithm A:

$3^{17} = 3 \times 3 = 129140163$,
so answer = 3.

Algorithm B (faster):

$3^{17} = (((3^2)^2)^2)^2 \times 3 = 129140163$, so answer = 3.

Algorithm C (fastest):

$(((((3^2 \bmod 10)^2 \bmod 10)^2 \bmod 10)^2 \bmod 10) \times 3 \bmod 10 = 3$.

Note that in Algorithm C, we never have to deal with integers larger than $81 = (n - 1)^2$. Thus it can be implemented using ordinary 64-bit integer arithmetic, even for very large n .

The BBP formula for π



In 1996, at the suggestion of Peter Borwein, Simon Plouffe used DHB's PSLQ integer relation program to discover this new formula for π :

$$\pi = \sum_{n=0}^{\infty} \frac{1}{16^n} \left(\frac{4}{8n+1} - \frac{2}{8n+4} - \frac{1}{8n+5} - \frac{1}{8n+6} \right)$$

This formula permits one to compute binary (or hexadecimal) digits of π beginning at an arbitrary starting position, using a very simple scheme that can run on any system with standard 64-bit or 128-bit arithmetic.

Recently it was proven that no base- b formulas of this type exist for π , except for when b is a power of two.

1. DHB, P. B. Borwein and S. Plouffe, "On the Rapid Computation of Various Polylogarithmic Constants," *Mathematics of Computation*, vol. 66, no. 218 (Apr 1997), pg. 903-913.
2. J. M. Borwein, W. F. Galway and D. Borwein, "Finding and Excluding b -ary Machin-Type BBP Formulae," *Canadian Journal of Mathematics*, vol. 56 (2004), pg 1339-1342.

Some other BBP-type formulas



$$\begin{aligned} \log \frac{9}{10} &= -\sum_{k=1}^{\infty} \frac{1}{k10^k} \\ \pi^2 &= \frac{9}{8} \sum_{k=0}^{\infty} \frac{1}{64^k} \left(\frac{16}{(6k+1)^2} - \frac{24}{(6k+2)^2} - \frac{8}{(6k+3)^2} - \frac{6}{(6k+4)^2} + \frac{1}{(6k+5)^2} \right) \\ \pi^2 &= \frac{2}{27} \sum_{k=0}^{\infty} \frac{1}{729^k} \left(\frac{243}{(12k+1)^2} - \frac{405}{(12k+2)^2} - \frac{81}{(12k+4)^2} - \frac{27}{(12k+5)^2} \right. \\ &\quad \left. - \frac{72}{(12k+6)^2} - \frac{9}{(12k+7)^2} - \frac{9}{(12k+8)^2} - \frac{5}{(12k+10)^2} + \frac{1}{(12k+11)^2} \right) \\ \zeta(3) &= \frac{1}{1792} \sum_{k=0}^{\infty} \frac{1}{2^{12k}} \left(\frac{6144}{(24k+1)^3} - \frac{43008}{(24k+2)^3} + \frac{24576}{(24k+3)^3} + \frac{30720}{(24k+4)^3} \right. \\ &\quad - \frac{1536}{(24k+5)^3} + \frac{3072}{(24k+6)^3} + \frac{768}{(24k+7)^3} - \frac{3072}{(24k+9)^3} - \frac{2688}{(24k+10)^3} \\ &\quad - \frac{192}{(24k+11)^3} - \frac{1536}{(24k+12)^3} - \frac{96}{(24k+13)^3} - \frac{672}{(24k+14)^3} - \frac{384}{(24k+15)^3} \\ &\quad + \frac{24}{(24k+17)^3} + \frac{48}{(24k+18)^3} - \frac{12}{(24k+19)^3} + \frac{120}{(24k+20)^3} + \frac{48}{(24k+21)^3} \\ &\quad \left. - \frac{42}{(24k+22)^3} + \frac{3}{(24k+23)^3} \right) \end{aligned}$$

BBP formulas and normality



Consider the general BBP-type constant

$$\alpha = \sum_{n=0}^{\infty} \frac{p(n)}{b^n q(n)}$$

where p and q are integer polynomials, $\deg p < \deg q$, and q has no zeroes for nonnegative arguments. Let $\{ \}$ denote fractional part.

In 2001, DHB and Richard Crandall proved that α is b -normal iff the sequence $x_0 = 0$, and

$$x_n = \left\{ bx_{n-1} + \frac{p(n)}{q(n)} \right\}$$

is equidistributed in the unit interval. Here “equidistributed” means that the sequence visits each subinterval $[c, d)$ with limiting frequency $d - c$.

DHB and R. E. Crandall, “On the Random Character of Fundamental Constant Expansions,” *Experimental Mathematics*, vol. 10, no. 2 (Jun 2001), pg. 175-190.

Two specific examples



Let $\{ \}$ denote fractional part, and consider the sequence $x_0 = 0$, and

$$x_n = \left\{ 2x_{n-1} + \frac{1}{n} \right\}$$

Then $\log 2$ is 2-normal iff this sequence is equidistributed in the unit interval.

Similarly, consider the sequence $x_0 = 0$, and

$$x_n = \left\{ 16x_{n-1} + \frac{120n^2 - 89n + 16}{512n^4 - 1024n^3 + 712n^2 - 206n + 21} \right\}$$

Then π is 16-normal (and hence 2-normal) iff this sequence is equidistributed in the unit interval.

A class of provably normal constants



DHB and Crandall have also shown that an infinite class of mathematical constants is 2-normal, including

$$\begin{aligned}\alpha_{2,3} &= \sum_{n=1}^{\infty} \frac{1}{3^n 2^{3^n}} \\ &= 0.041883680831502985071252898624571682426096 \dots_{10} \\ &= 0.0ab8e38f684bda12f684bf35ba781948b0fcd6e9e0 \dots_{16}\end{aligned}$$

This constant was proven 2-normal by Stoneham in 1971, but we have extended this to the case where (2,3) are any pair (p,q) of relatively prime integers > 1. We also extended this result to an uncountable class:

$$\alpha_{2,3}(r) = \sum_{n=1}^{\infty} \frac{1}{3^n 2^{3^n + r_n}}$$

Here r_n is the n -th bit in the binary expansion of r in (0,1). These constants are all distinct.

DHB and R. E. Crandall, "Random Generators and Normal Numbers," *Experimental Mathematics*, vol. 11, no. 4 (2002), pg. 527-546.

A “hot spot” lemma



Given the real constant α , if there exists some B such that for every subinterval $[c, d)$ of $[0, 1)$,

$$\limsup_{m \geq 1} \frac{\#\{0 \leq j < m \mid \{b^j \alpha\} \in [c, d)\}}{m(d - c)} \leq B$$

then α is b -normal.

In other words, if α is not b -normal, then there is some interval $[c, d)$ that is visited 10 times too often by shifts of the base- b expansion of α ; there is some other interval $[c', d')$ that is visited 100 times too often; there is some other interval $[c'', d'')$ that is visited 1000 times too often, etc. However, one cannot conclude that these intervals are nested.

L. Kuipers and H. Niederreiter, *Uniform Distribution of Sequences*, 1974, pg. 77.

A strong “hot spot” lemma



Recently DHB and Michal Misiurewicz proved a stronger version of this result, using methods of ergodic theory:

Let $0.x_1x_2\dots x_n$ be the base- b expansion of x to position n . If for every x in $(0,1)$,

$$\liminf_{n \geq 1} \limsup_{m \geq 1} \frac{\#\{0 \leq j < m \mid \{b^j \alpha\} \in [0.x_1x_2\dots x_n, 0.x_1x_2\dots x_n + b^{-n})\}}{mb^{-n}} < \infty$$

then α is b -normal.

In other words, if α is not b -normal, then there is at least one x in $(0,1)$ such that shifts of the base- b expansion of α visit all sufficiently small digit neighborhoods of x too often, by an arbitrarily large factor.

DHB and M. Misiurewicz, “A Strong Hot Spot Theorem,” *Proceedings of the American Mathematical Society*, vol. 134 (2006), no. 9, pg. 2495-2501.

The BBP sequence corresponding to $\alpha_{2,3}$



It is fairly easy to show that the BBP sequence (z_n) corresponding to

$$\alpha_{2,3} = \sum_{n=1}^{\infty} \frac{1}{3^n 2^{3^n}}$$

is the following. Note that each section is repeated three times, and the sequence evenly fills in the unit interval with all fractions of the form $k 3^{-p}$.

$$\begin{aligned}
 &0, 0, 0, \\
 &\frac{1}{3}, \frac{2}{3}, \frac{1}{3}, \frac{2}{3}, \frac{1}{3}, \frac{2}{3}, \\
 &\frac{4}{9}, \frac{8}{9}, \frac{7}{9}, \frac{5}{9}, \frac{1}{9}, \frac{2}{9}, \quad (\text{repeated 3 times}), \\
 &\frac{13}{27}, \frac{26}{27}, \frac{25}{27}, \frac{23}{27}, \frac{19}{27}, \frac{11}{27}, \frac{22}{27}, \frac{17}{27}, \frac{7}{27}, \frac{14}{27}, \frac{1}{27}, \frac{2}{27}, \frac{4}{27}, \frac{8}{27}, \frac{16}{27}, \frac{5}{27}, \\
 &\frac{10}{27}, \frac{20}{27}, \quad (\text{repeated 3 times}), \text{ etc.}
 \end{aligned}$$

It is also easy to show that the sequence of shifted bits satisfies, for all $n > 0$,

$$|\{2^n \alpha_{2,3}\} - z_n| < \frac{1}{2n}$$

Proof that $\alpha_{2,3}$ is 2-normal



Given any x in $(0,1)$, let $c = 0.x_1x_2\dots x_n$, and $d = 0.x_1x_2\dots(x_n+1)$. Let m be any integer $> 2^{2n}$, and let 3^p such that $3^p \leq m < 3^{p+1}$. Note that for $j > 2^n$, $[c - 1/(2^j), d + 1/(2^j))$ is a subset of $[c - 2^{-n-1}, d + 2^{-n-1})$. Since the length of this interval is 2^{-n+1} , it contains at most $3^p 2^{-n+1} + 1$ instances of $k 3^{-p}$.

Therefore

$$\begin{aligned} \frac{\#\{0 \leq j < m \mid \{2^j \alpha\} \in [c, d)\}}{m 2^{-n}} &\leq \frac{2^n + \#\{2^n \leq j < m \mid \{z_j\} \in [c - 2^{-n-1}, d + 2^{-n-1})\}}{m 2^{-n}} \\ &\leq \frac{2^n + 3(3^p 2^{-n+1} + 1)}{m 2^{-n}} < 8 \end{aligned}$$

Thus by the hot spot lemma, $\alpha_{2,3}$ is 2-normal.

See paper by DHB and Misiurewicz for full details.

$\alpha_{2,3}$ is not 6-normal



It is also possible to establish non-normality results for $\alpha_{2,3}$ in certain number bases, such as base 6. Note that we can write

$$\{6^n \alpha_{2,3}\} = \left\{ \sum_{m=1}^{\lfloor \log_3 n \rfloor} 3^{n-m} 2^{n-3^m} \right\} + \left\{ \sum_{m=\lfloor \log_3 n \rfloor + 1}^{\infty} 3^{n-m} 2^{n-3^m} \right\}.$$

The first portion of this expression is zero, since all of the terms in the summation are zero. When $n = 3^m$, the second portion is, very accurately,

$$\{6^{3^m} \alpha_{2,3}\} \approx \frac{\left(\frac{3}{4}\right)^{3^m}}{3^{m+1}}$$

Thus the base-6 expansion of $\alpha_{2,3}$ has long stretches of zeroes beginning at positions $3^m + 1$. This observation can be fashioned into a rigorous proof of non-normality.

DHB, "A Non-Normality Result," manuscript, Aug 2007, <http://crd.lbl.gov/~dhbailey/dhbpapers/alpha-6.pdf>.

A pseudorandom number generator based on the binary digits of $\alpha_{2,3}$



Given a seed q in the range $3^{33} + 100 < q < 2^{53}$, use the first line to compute x_0 , and use the second line to compute all successive iterates:

$$x_0 = (2^{q-3^{33}} \cdot \lfloor 3^{33}/2 \rfloor) \bmod 3^{33}$$

$$x_k = (2^{53} \cdot x_{k-1}) \bmod 3^{33}$$

Divide the results by 3^{33} to obtain pseudorandom 64-bit floating-point iterates in $(0,1)$.

This generator has several desirable properties:

- Iterates contain successive 53-bit sections of the binary digits of $\alpha_{2,3}$.
- It is not subject to power-of-two stride problems that plague other schemes.
- It passes all standard tests for randomness.
- It is well-suited for parallel processing – each individual processor can quickly jump to its own starting point in the sequence.
- Efficient implementations are as fast as several other widely used schemes.

DHB, "A Pseudo-Random Number Generator Based on Normal Numbers," Dec 2004,
<http://crd.lbl.gov/~dhbailey/dhbpapers/normal-random.pdf>.